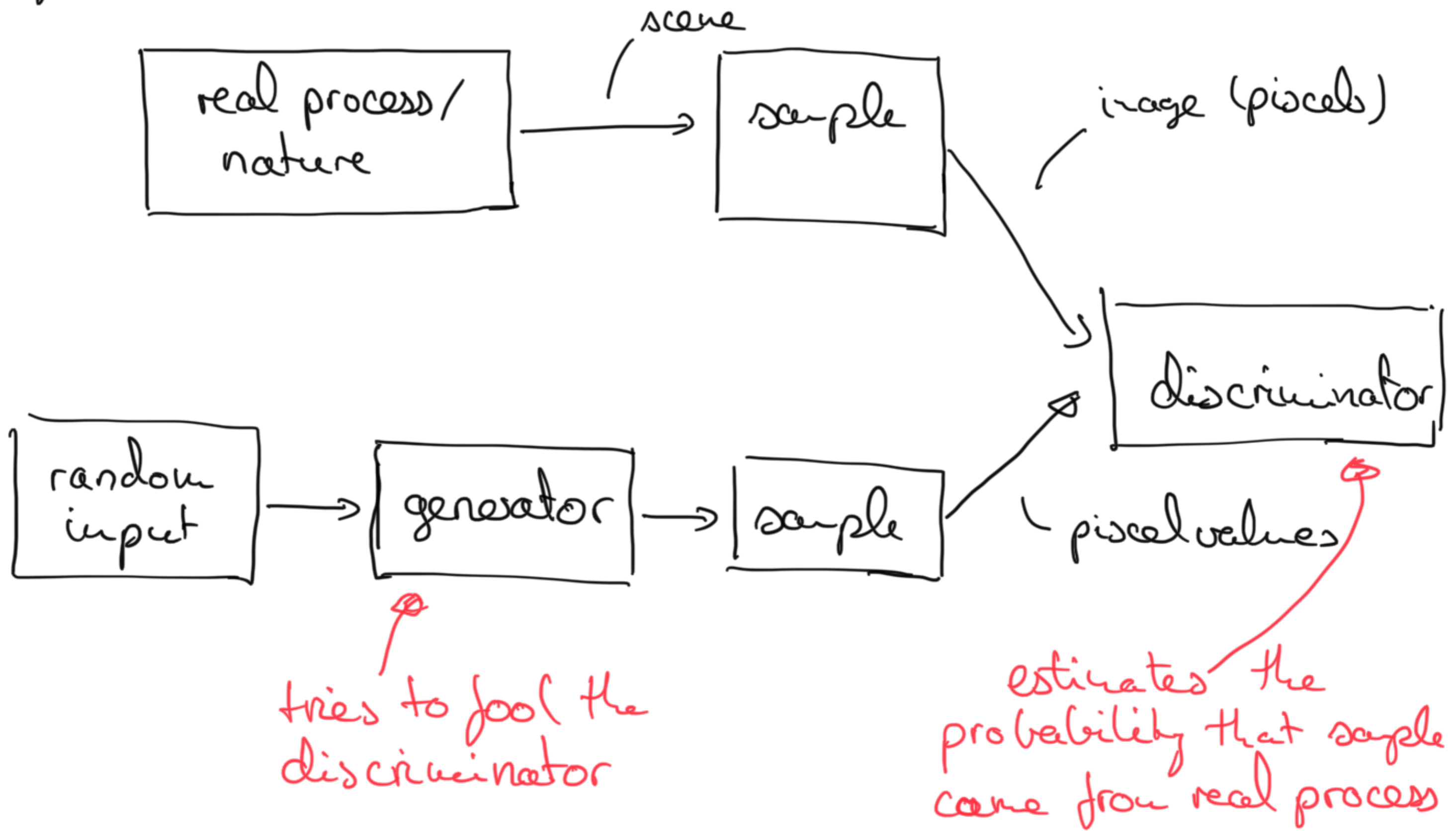


Brief overview of game theory

- Literature :
- A. R. Katin, J. Peseo, "Game Theory, Alive", AMS, 2016
 - N. Cesa Bianchi, G. Lugosi, "Predictions, Learning and Games", Cambridge University Press, 2006

1. Motivation

- i). Learning system is interacting with an environment
→ for example in the smart grid, different entities make decisions about production, consumption, and storage of energy and have conflicting interests.
- ii). generative adversarial networks:



iii). convex optimization : \rightarrow design an algorithm that works for any function in a given class
 \rightarrow what is the worst-case rate?

2. Two-player zero-sum games

- player I can choose between actions $1 \dots m$
- player II can choose between actions $1 \dots n$

→ player I receives a_{ij}

→ player II receives $-a_{ij}$

for choosing action i (player I) and action j (player II)

- a_{ij} , $i=1, \dots, m$, $j=1, \dots, n$ is the payoff

◦ Example

		P. II	
		①	②
P. I	①	1	2
	②	3	4

→ How should they play → think about the worst-case.

→ p. I has an incentive to play ②

→ p. II has an incentive to play ①

→ this leads to an eq. → no player has an incentive to deviate from his worst-case strategy

		p. II	
		①	②
p. I	①	1	0.5
	②	0	2

→ p. I has an incentive for ①

→ p. II has an incentive for ①

∴ does not lead to an eq.

→ this one

→ How can we resolve this (from a mathematical point of view)?

→ p. I chooses action i with prob. x_i

→ p. II chooses action j with prob. y_j

→ $x = (x_1, \dots, x_m) \in \Delta_m$

→ $y = (y_1, \dots, y_n) \in \Delta_n$

• gain for p. I: $x^T A y$

• gain for p. II: $-x^T A y$

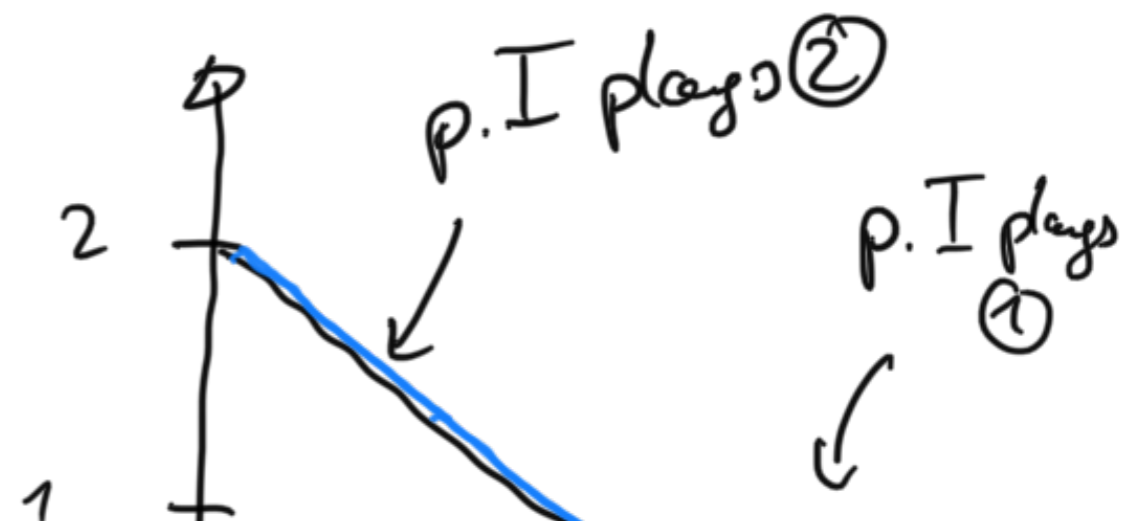
from p. I's perspective

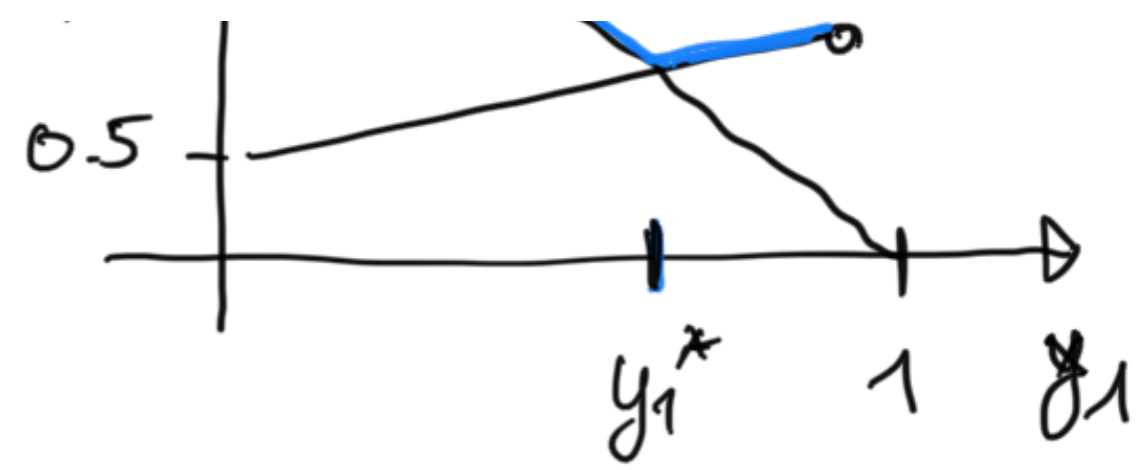
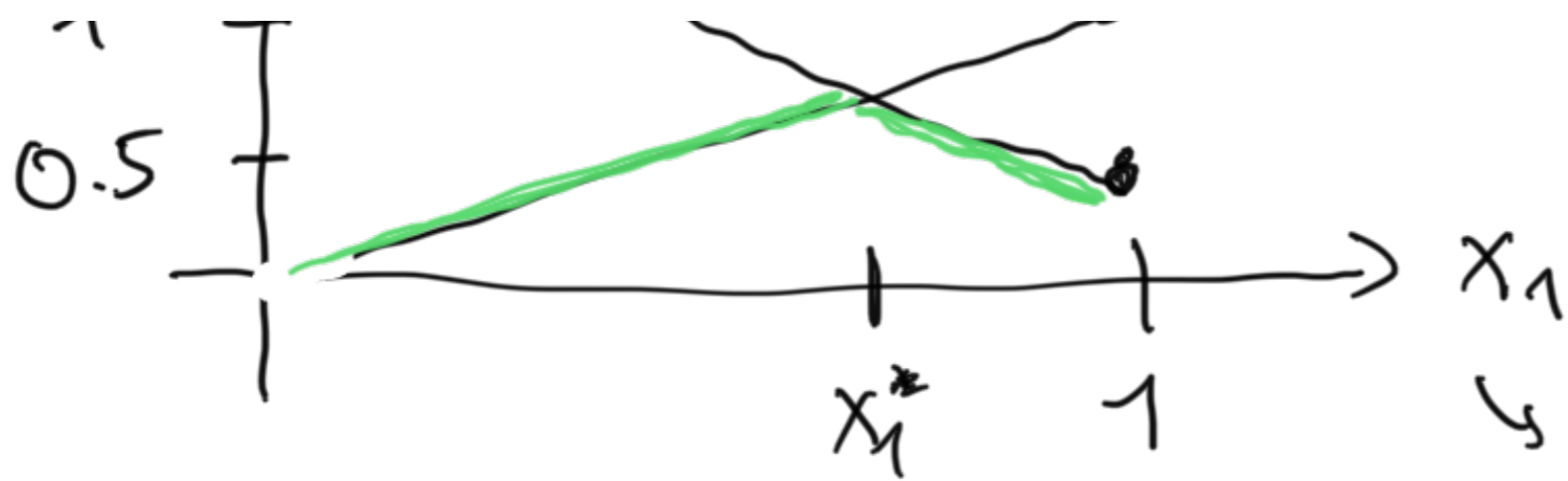


p. II plays ②

p. II plays ①

from p. II's persp.





$$\downarrow \\ x_2 = 1 - x_1$$

$$x_1^* = 2(1 - x_1^*) + 0.5x_1^* \quad \leadsto \quad \begin{aligned} x_1^* &= 4/5 \\ x_2^* &= 1/5 \end{aligned}$$

$$y_1^* = 2/3, \quad y_2^* = 1/3$$

- Assume p.I plays (x_1^*, x_2^*) , the payoff for p.II is $-4/5$ for action ① and $-4/5$ for action ②
 \rightarrow p.II has no incentive in deviating

- Def. (x^*, y^*) is a Nash eq. or equilibrium if

$$x^{*T} A y^* \leq x^{*T} A y \quad \forall y \in \Delta_n$$

$$x^T A y \geq x^T A y^* \quad \forall x \in \Delta_m$$

Prop. The following are equivalent:

(i) there exists a Nash eq.

$$(ii) \quad v = \max_{x \in \Delta_m} \underbrace{\min_{y \in \Delta_n} x^T A y}_{\text{worst case from p. I's persp.}} = \min_{y \in \Delta_n} \underbrace{\max_{x \in \Delta_m} x^T A y}_{\text{worst case from p. II's persp.}}$$

Remarks: • interpretation of (ii)
 → can compute strategies to achieve v

$$\max_{x \in \Delta_m} \underbrace{\min_{y \in \Delta_n} x^T A y}_{\text{p. II knows about}} = \min_{y \in \Delta_n} \underbrace{\max_{x \in \Delta_m} x^T A y}_{\text{p. I knows about}}$$

$$\sigma_1 = \min_{y \in \Delta_n} \max_{x \in \Delta_m} x^T A y \geq \max_{x \in \Delta_m} \min_{y \in \Delta_n} x^T A y = \sigma_2$$

• Proof (i) \Rightarrow (ii) (x^*, y^*) is a Nash eq.

$$x^{*T} A y^* = \underbrace{\max_{x \in \Delta_m} x^T A y^*}_{g(y^*)} \geq \min_{y \in \Delta_n} \max_{x \in \Delta_m} x^T A y = \sigma_1$$

$$x^{*T} A y^* = \underbrace{\min_{y \in \Delta_n} x^{*T} A y}_{\tilde{g}(x^*)} \leq \max_{x \in \Delta_m} \min_{y \in \Delta_n} x^T A y = \sigma_2$$

$$\Rightarrow \sigma_1 = \sigma_2$$

$$(ii) \Rightarrow (i) \quad x^* = \operatorname{argmax}_{x \in \Delta_m} \min_{y \in \Delta_n} x^T A y$$

$$y^* = \operatorname{arg\,min}_{y \in \Delta_n} \max_{x \in \Delta_n} x^T A y$$

Theorem (Nash) There exists a Nash eq.
(holds also for general sum games)

Proof. Idea: introduce a map

$$T: \Delta_n \times \Delta_n \rightarrow \Delta_n \times \Delta_n$$

$$T(x, y) = (x', y'), \text{ such that}$$

(x', y') improves over (x, y) . previous
reward
↓

$$\rightarrow \text{define } c_i(x, y) = \max \left\{ \sum_{j=1}^n a_{ij} y_j - x^T A y, 0 \right\}$$

$$d_j(x, y) = \max \left\{ x^T A y - \sum_{i=1}^n a_{ij} x_i, 0 \right\}$$

$$x' + r: (x, y)$$

$$x_i' = \frac{x_i + d_i}{1 + \sum_{k=1}^m c_k(x, y)} \rightarrow \text{normalization}$$

$$y_j' = \frac{y_j + d_j}{1 + \sum_{k=1}^n d_k(x, y)} \rightarrow \text{norm.}$$

- T is continuous, it maps a compact convex set to itself.
- Brouwer's theorem to conclude that $\exists (x^*, y^*) : T(x^*, y^*) = (x^*, y^*)$.
- (x^*, y^*) corresponds to a Nash eq.

→ Alternative decision making

2. Market Prediction

• Stock market prediction:

→ different algorithms make predictions about buying/selling stocks

→ at each day we follow one of the alg. predictions

→ at the end of the day we observe the outcome of all the different algorithms including the one we decided to follow

→ 1500 days horizon → which algorithm should you follow to maximize revenue?

• Formalization:

... .. between

• A decider has to pick one of actions $\{1, \dots, n\}$

• At each round t :

(i) for each action $i \in \{1, \dots, n\}$ the environment generates a loss $l_i^t \in [0, 1]$.

(ii) the decider chooses a prob. distribution $y^t \in \Delta_n$ over actions based on $l^1, l^2, \dots, l^{t-1}, y^1, y^2, \dots, y^{t-1}$, where $l^j = (l_1^j, l_2^j, \dots, l_n^j)$

(iii) the decider suffers the loss $l^{tT} y^t$.

o Performance measure:

$$R_{T_r}(l, D) = \sum_{t=1}^{T_r} y^{+T} l^t - T_r \min_{y \in \Delta_n} \left(\frac{1}{T_r} \sum_{t=1}^{T_r} l^t \right) y$$

empirical
average
over the
observed losses

$$= \sum_{t=1}^{T_r} y^{+T} l^t - \min_{i \in \{1, \dots, n\}} \sum_{k=1}^{T_r} l_i^k$$

→ Strategy "multiplicative weights"

$$\hookrightarrow R_{T_r}(l, D) \leq \sqrt{\frac{\log(n) T_r}{2}}$$